

# Proof Theory of Classical Logic

Its Basics with an Emphasis on Quantitative Aspects

Short course at Notre Dame

Feb 5: Pudlak's Exponential Arithmetic

# Outline

Pudlak's exponential arithmetic PEX

Analysing unprovability and provability in PEX

# Addenda

## Upper bound for midsequent theorem

A proof of depth  $n$  has at most  $2^{n-1}$  quantifier inferences (we temporarily define  $2^{-1} = 0$ ).

So, a regular cut-free proof of depth  $n$ , whose final sequent consists of prenex formulas only, can be converted to a midsequent proof of depth  $n + 2^{n-1}$ . So of depth  $2^n$ .

So, a proof  $\mathcal{P}$  of regular prenex sequent can be converted to a midsequent proof having depth  $\leq 2_{r(\mathcal{P})+1}^{d(\mathcal{P})}$ .

## Tautologies in predicate logic

A predicate formula is a **tautology** if it can be obtained from a propositional tautology by substituting predicate formulas for its atoms (same formulas for same atoms).

## Numerals

Having a language containing 0 and S, the term  $S(S(\dots S(0)\dots))$  with  $n$  occurrences of S is denoted  $\bar{n}$  and called ( **$n$ -th numeral**). Thus 0 and  $\bar{0}$  are the same terms.

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# Pudlak's exponential arithmetic PEX

Axioms about 0, S, and +

$$\text{Q1: } \forall x \forall y (S(x) = S(y) \rightarrow x = y),$$

$$\text{Q2: } \forall x (S(x) \neq 0),$$

$$\text{Q4: } \forall x (x + 0 = x),$$

$$\text{Q5: } \forall x \forall y (x + S(y) = S(x + y)),$$

$$\forall x \forall y (x + y = y + x),$$

$$\forall x \forall y \forall z (x + (y + z) = (x + y) + z),$$

Axioms about E

$$E(0) = S(0),$$

$$\forall x (E(S(x)) = E(x) + E(x)),$$

Axioms about P

$$P(0),$$

$$\forall x (P(x) \rightarrow P(S(x))).$$

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# Motivating questions, models

## Questions

Can  $PEX \vdash \forall x P(x)$ ? Can  $PEX \vdash P(E^{(60)}(0))$ ?

In our setting, these are questions about provability of sequents  $\langle PEX \Rightarrow \forall x P(x) \rangle$  and  $\langle PEX \Rightarrow P(E^{(60)}(0)) \rangle$ , where  $PEX$  is a set containing the 10 axioms of the theory  $PEX$  plus 7 identity axioms (listed on next frame).

## Some models of $PEX$

1.  $M_0 = \langle \mathbb{N}, 0^{\mathbb{N}}, (a \mapsto a + 1)^{\mathbb{N}}, +^{\mathbb{N}}, (a \mapsto 2^a)^{\mathbb{N}}, \mathbb{N} \rangle$ .
2. Let  $M = \langle M, 0^M, S^M, +^M, \cdot^M \rangle$  be a non-standard model of PA. Take  $\langle M, 0^M, S^M, +^M, (a \mapsto 2^a)^M, \mathbb{N} \rangle$ .
3. In the same model, change the last component as follows. Fix  $a_0 \in M$ , and send  $P$  to the set  $a_0 + \mathbb{N} = \{ a ; \exists n \in \mathbb{N} (a \leq^M a_0 +^M n) \}$ .
4.  $\langle \mathbb{R} \cap [0, \infty), 0^{\mathbb{R}}, (a \mapsto a + 1)^{\mathbb{R}}, +^{\mathbb{R}}, (a \mapsto 2^a)^{\mathbb{R}}, \mathbb{N} \rangle$ .

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# Identity axioms

## Seven identity axioms in the set PEX

$$\forall x(x = x),$$

$$\forall x\forall y(x = y \rightarrow y = x),$$

$$\forall x\forall y\forall z(x = y \rightarrow (y = z \rightarrow x = z)),$$

$$\forall x\forall y(x = y \rightarrow (S(x) = S(y))),$$

$$\forall x_1\forall x_2\forall y_1\forall y_2(x_1 = y_1 \ \& \ x_2 = y_2 \rightarrow (x_1 + x_2 = y_1 + y_2)),$$

$$\forall x\forall y(x = y \rightarrow (E(x) = E(y))),$$

$$\forall x\forall y(x = y \rightarrow (P(x) \rightarrow P(y))).$$

## Example proof

Constructing a proof of  $\langle \text{PEX} \Rightarrow P(E^{(60)}(0)) \rangle$

1. Find number  $m$  such that  $\bar{m} = E^{(60)}(0)$ .
2. Prove the sequent  $\langle \text{PEX} \Rightarrow \bar{m} = E^{(60)}(0) \rangle$ .
3. Prove the sequent  
 $\langle \text{PEX}, P(0), P(\bar{0}) \rightarrow P(\bar{1}), \dots, P(\overline{m-1}) \rightarrow P(\bar{m}) \Rightarrow P(\bar{m}) \rangle$ .
4. Un-substitute  $0, \dots, \overline{m-1}$ , and obtain  $\langle \text{PEX} \Rightarrow P(\bar{m}) \rangle$ .
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6. Un-substitute, i.e. use  $\forall I$  twice, to obtain  
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7. Cuts on sequents from 2 and 4 yield  $\langle \text{PEX} \Rightarrow P(E^{(60)}(0)) \rangle$ .



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